

Parallel Functional Programming with Interaction Nets

Marc Thatcher

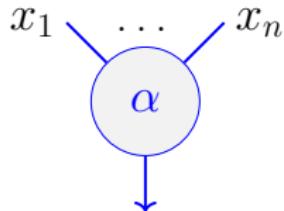
`m.thatcher@sussex.ac.uk`

Department of Informatics, University of Sussex

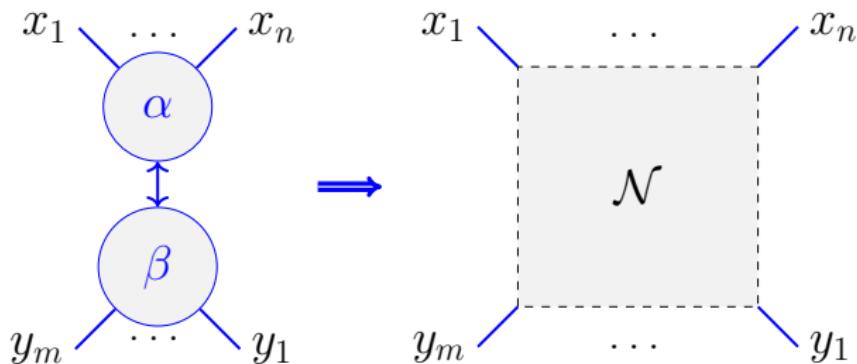
40th British Colloquium for Theoretical Computer Science
University of Bath, April 4th–5th 2024

Interaction nets (Lafont,1990)

Finite set of *user-defined* agents:

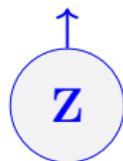


Finite set of *user-defined* rewrite rules:

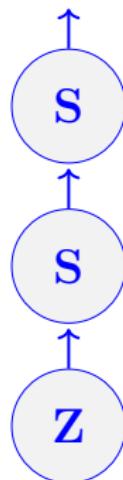


Example - Unary numbers

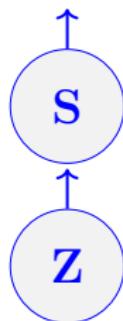
Zero



Two

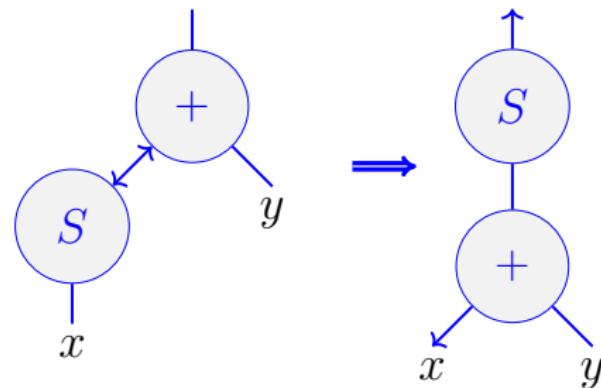
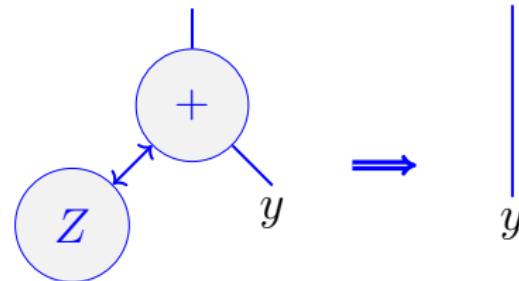


One

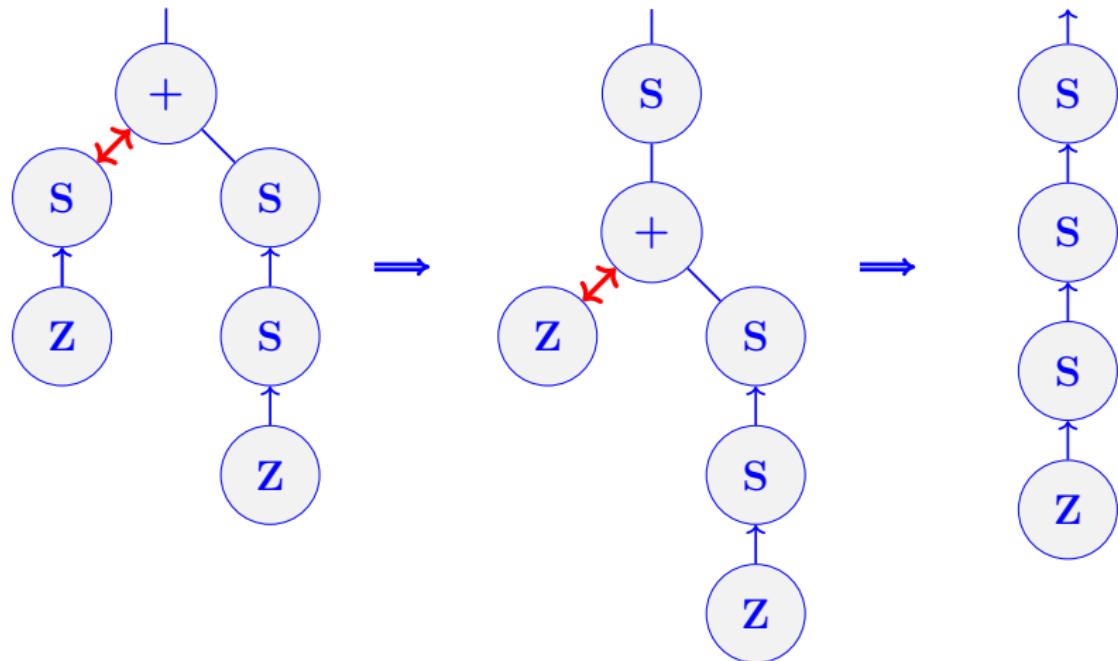


etc. ...

Example - Unary number addition



Example function - Unary number addition

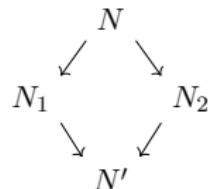


Properties as a programming language

- ▶ Turing complete
- ▶ Pattern matching
- ▶ Constant time rewrites
- ▶ Visual debugging

- ▶ Local reductions ; Shared computations

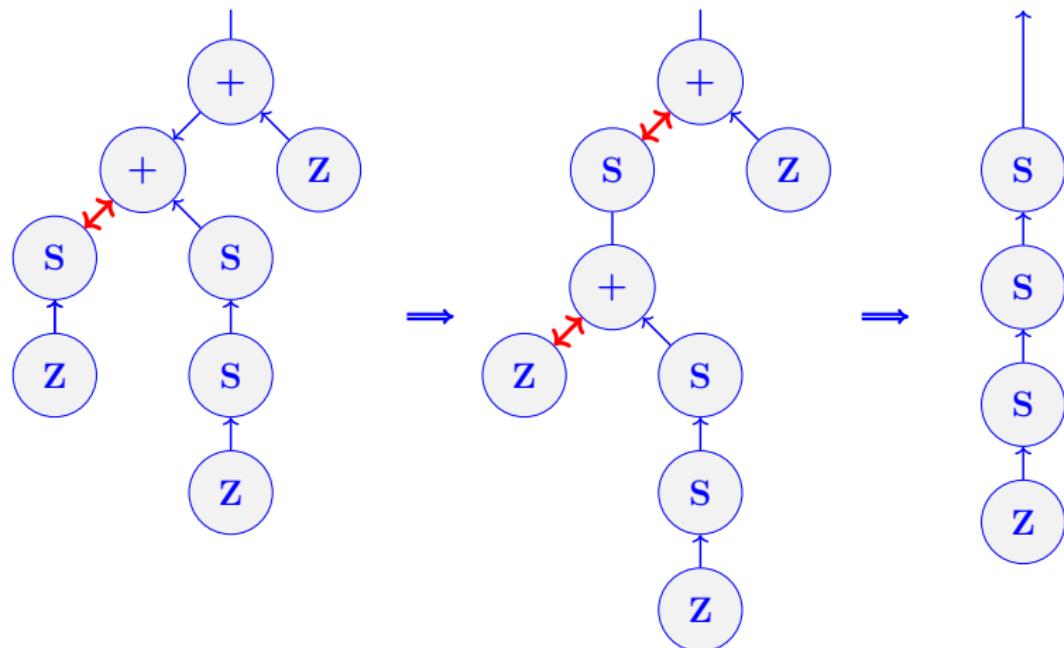
- ▶ One-step confluence



- ▶ Explicit mandatory memory management – **no GC!**

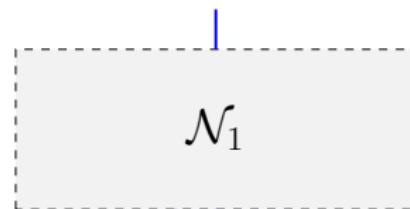
→ Natural parallel execution

Parallel evaluation



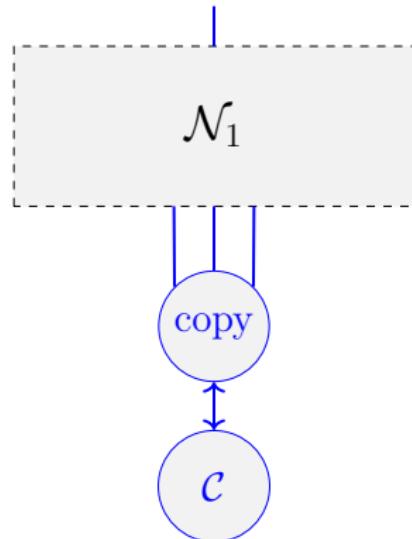
Sharing

Assume $\mathcal{N}_2 \rightarrow^* \mathcal{C}$

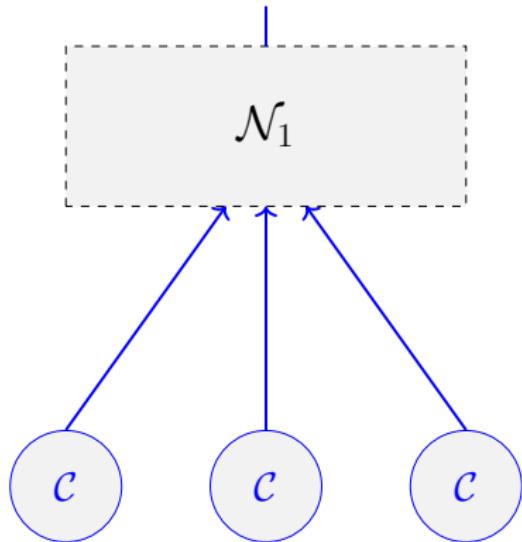


\mathcal{N}_2

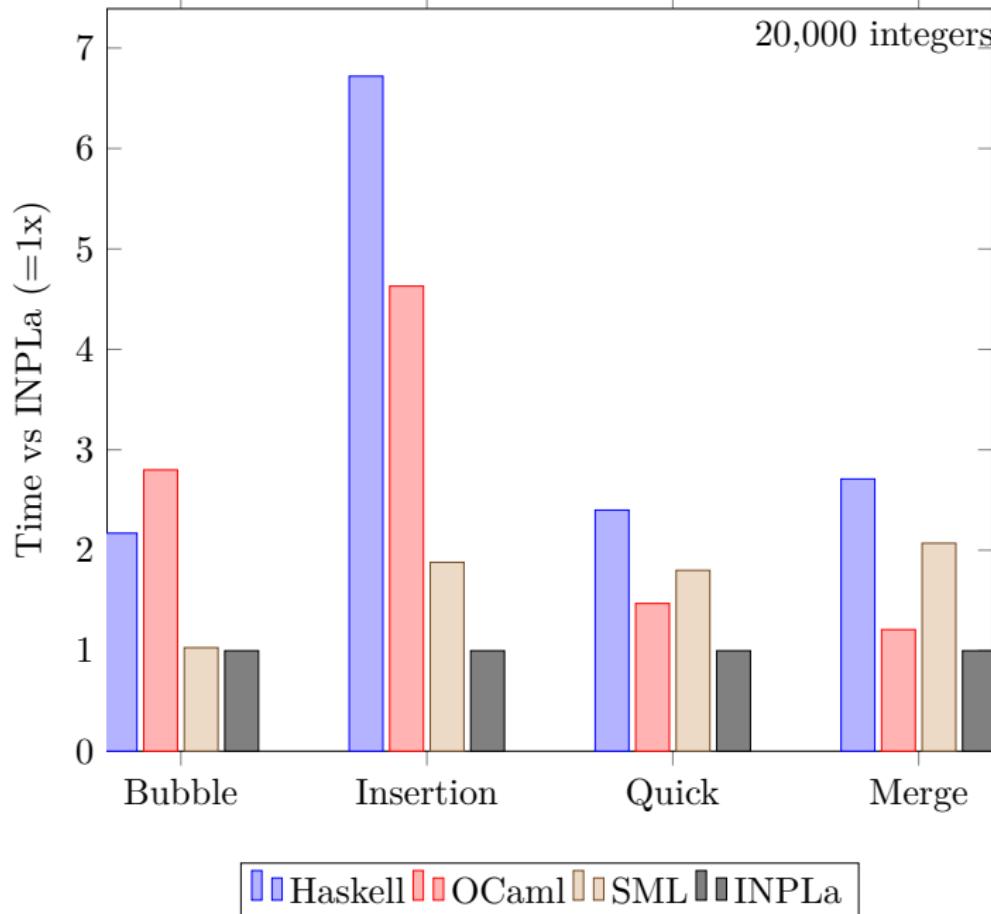
Sharing



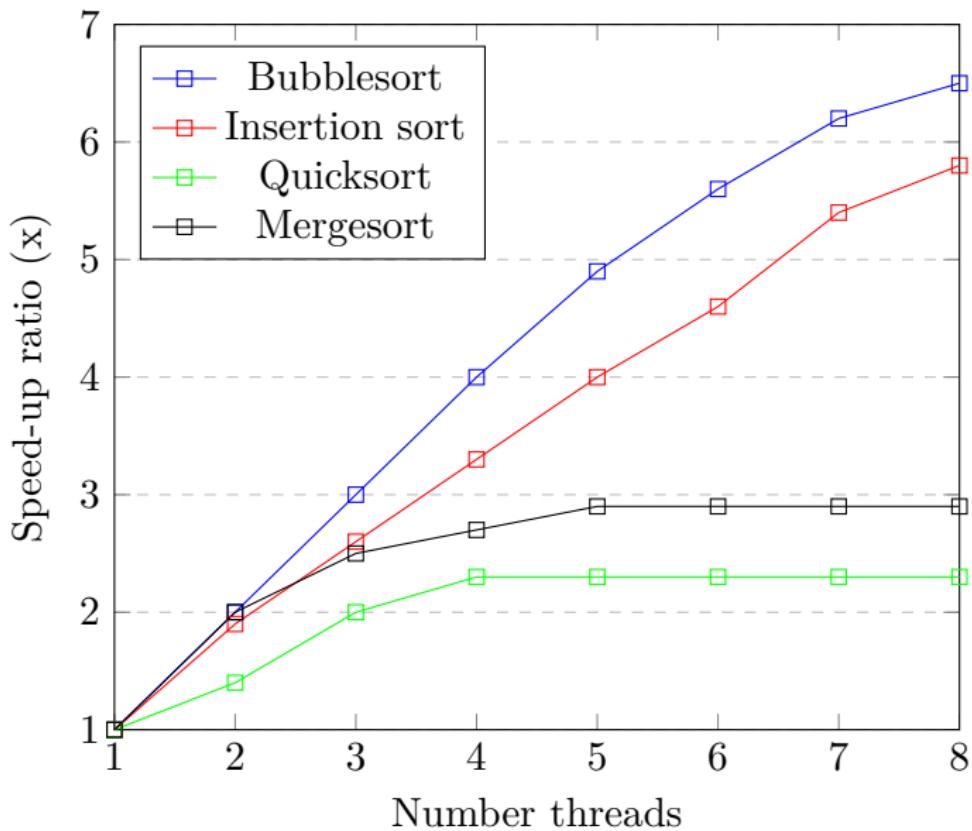
Sharing



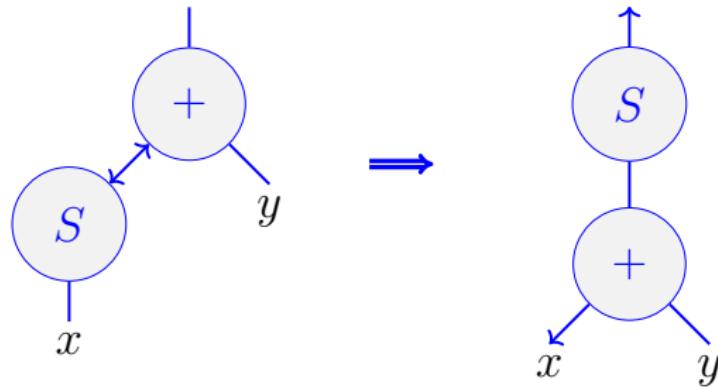
Impact of parallelism - benchmark results



Impact of parallelism - benchmark results



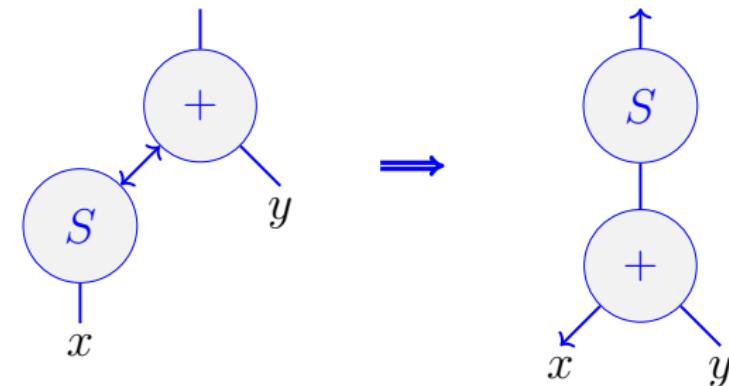
Towards a programming language v.1¹



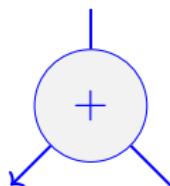
`add(result,y)><S(x) => result~S(aux), add(aux,y)~x`

¹Sato, 2014 ; <https://github.com/inpla/inpla>

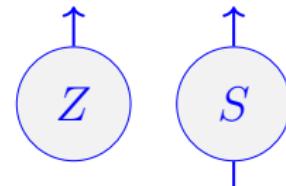
Towards a programming language v.2



Functions:

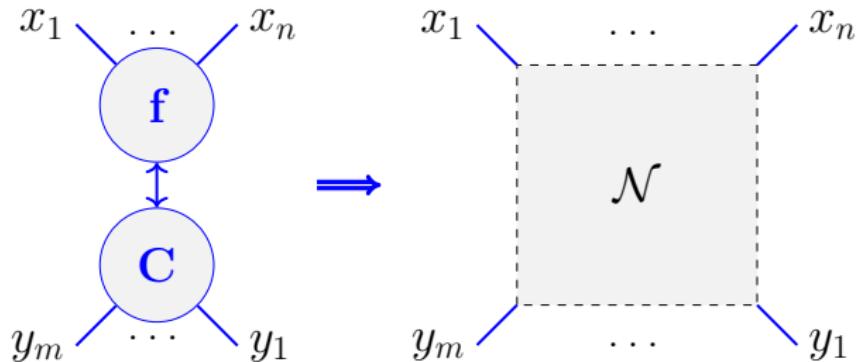


Constructors:



FLIN - a Functional Language for Interaction Nets²

If f is a function and C is a constructor:



then:

$$f(C(\vec{y}), \vec{x}') = N(\vec{x}', \vec{y})$$

where:

$$N = f \dots | C \dots | \vec{y} | \dots$$

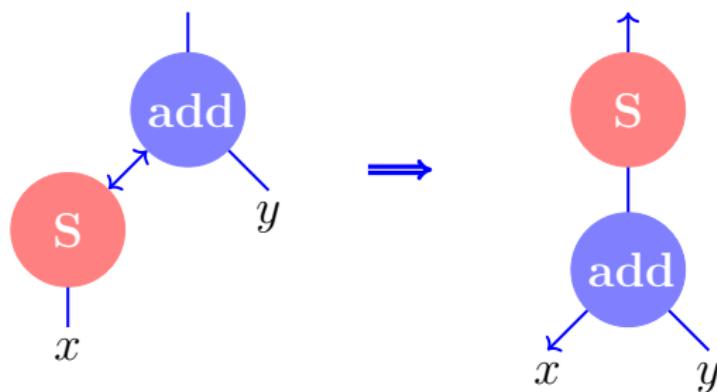
and

$$\vec{x}' = \vec{x} \text{ adjusted for output.}$$

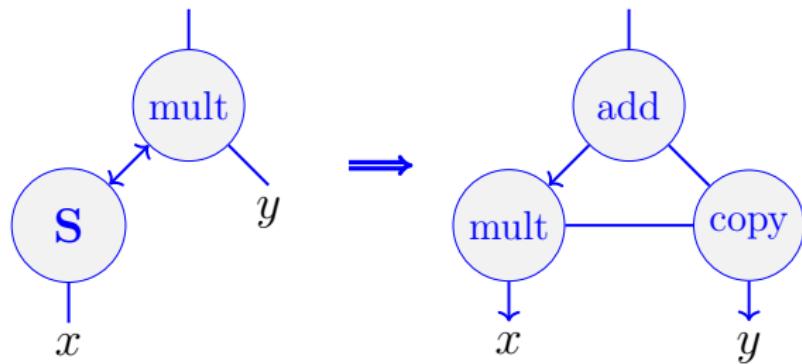
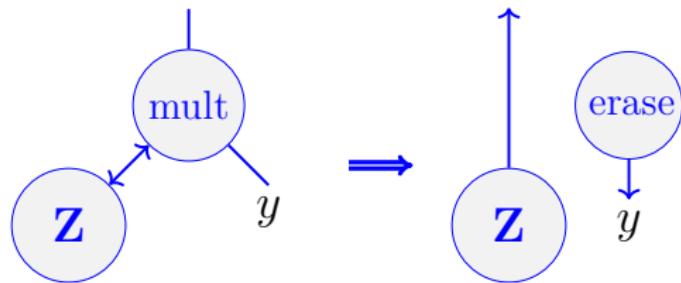
²<https://github.com/inpla/train>

FLIN \cong Interaction Nets

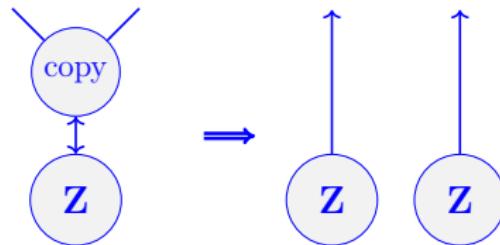
$$\text{add } S(x, y) = S(\text{add}(x, y))$$



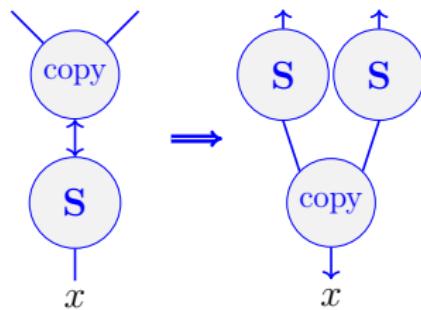
Non-functional interaction net rules



FLIN syntax for non-functions

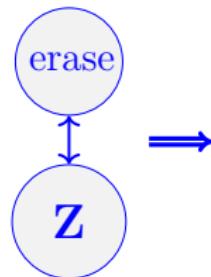


$\text{copy } Z = (Z, Z)$

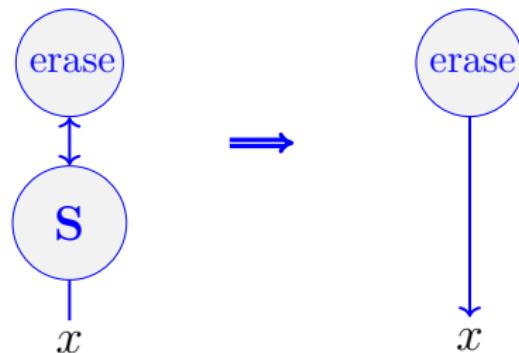


$\text{copy } S(x) = \text{let } (x_1, x_2) = \text{copy } x \text{ in } (S(x_1), S(x_2))$

FLIN syntax for non-functions



{erase><Z => }



{erase><S(x) => erase~x}

FLIN examples

```
mult (Z,y)      = (Z,{erase~y})  
mult (S(x),y) = let (y1,y2)=dup y in  
                  add (y1,(mult (x,y2)))
```

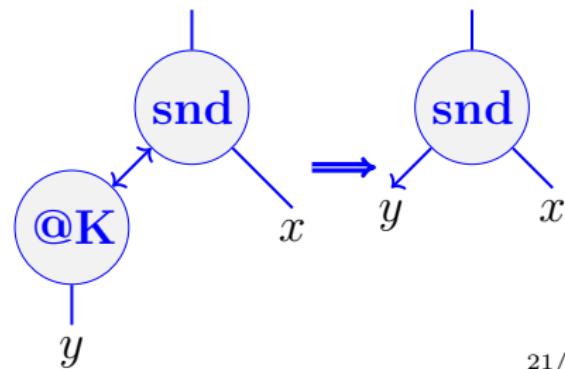
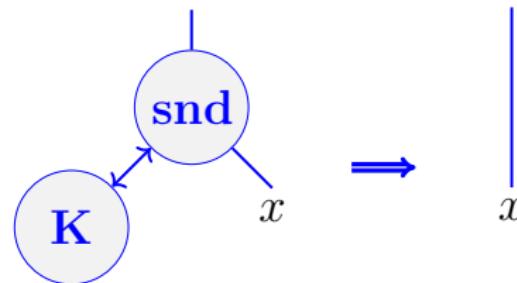
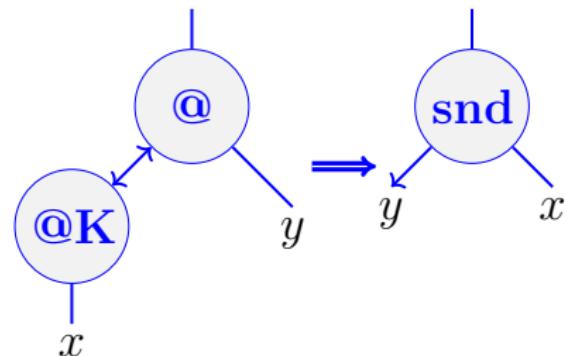
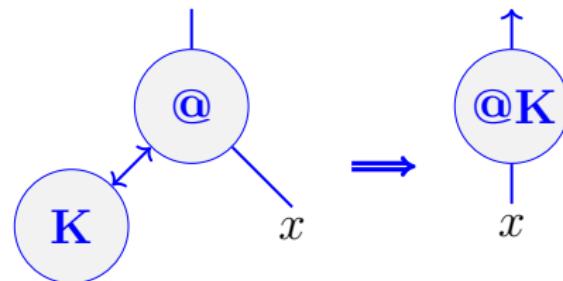
```
mult' (Z,y)      = snd(y,Z)  
snd   (Z,x)      = x  
snd   (S(y),x) = snd(y,x)
```

```
fib Z      = Z  
fib S(x)  = fibS x  
fibS Z     = S(Z)  
fibS S(x) = let (x1,x2)=dup x in  
                  add ((fibS x1),(fib x2))
```

```
append ([] ,ys)      = ys  
append ((x:xs),ys) = x:(append (xs,ys))
```

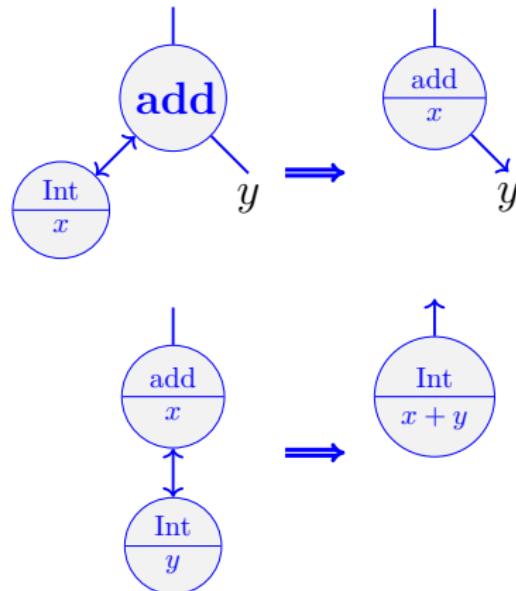
Computational power of Functional Interaction Nets

$$\mathbf{K}xy = x$$



Extension - Attributes

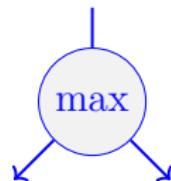
Hold values within agents - ints, bools, strings etc. & tuples of.
(Fernández, Mackie, Pinto 2001)



cf. λ -calculus \rightarrow PCF.

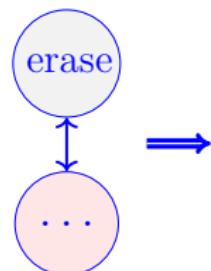
Extensions – further work

Multiple principal ports



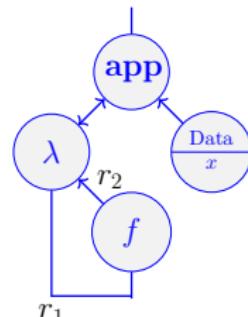
INMPP (Alexiev, 1999);
Macros (Sinot, Mackie 2005)

Generic rules

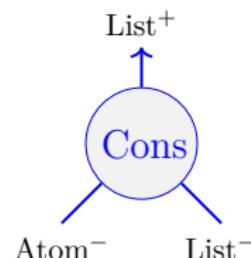


(Jiresch, 2012)

Higher order functions



Type system



(Lafont, 1990); (Fernández, 1998)

Conclusions

- ▶ Interaction nets: asynchronous parallel computation.
- ▶ INPLa implementation has encouraging benchmarks.
github.com/inpla/inpla
- ▶ FLIN - function-constructor language maps 1:1 to interaction nets.
- ▶ FLIN → INPLa transpiler.
github.com/inpla/train
- ▶ FLIN – programming or intermediate language for more complete language.

Lots more to do!